

[HMU]

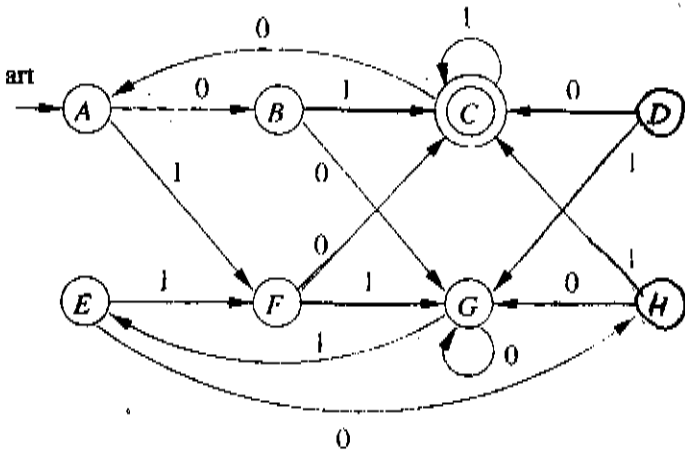


Figure 4.8: An automaton with equivalent states

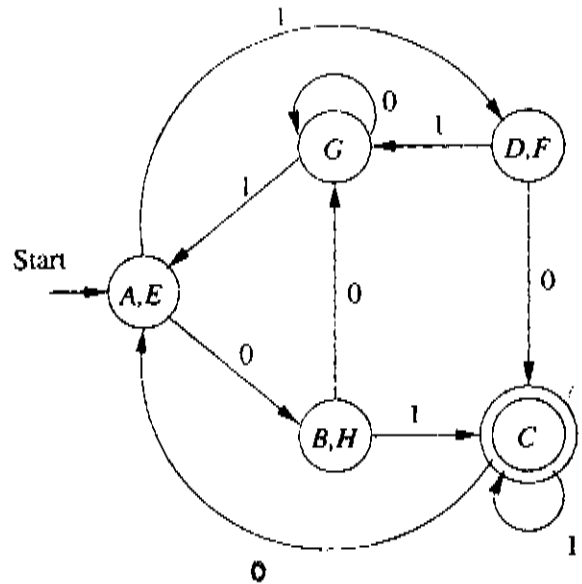


Figure 4.12: Minimum-state DFA equivalent to Fig. 4.8

**Example 4.19:** Let us execute the table-filling algorithm on the DFA of Fig 4.8. The final table is shown in Fig. 4.9, where an  $x$  indicates pairs of distinguishable states, and the blank squares indicate those pairs that have been found equivalent. Initially, there are no  $x$ 's in the table.

B	x						
C	x	x					
D	x	x	x				
E		x	x	x			
F	x	x	x		x		
G	x	x	x	x	x	x	
H	x		x	x	x	x	x
	A	B	C	D	E	F	G

Figure 4.9: Table of state inequivalences

For the basis, since  $C$  is the only accepting state, we put  $x$ 's in each pair that involves  $C$ . Now that we know some distinguishable pairs, we can discover others. For instance, since  $\{C, H\}$  is distinguishable, and states  $E$  and  $F$  go to  $H$  and  $C$ , respectively, on input 0, we know that  $\{E, F\}$  is also a distinguishable pair. In fact, all the  $x$ 's in Fig. 4.9 with the exception of the pair  $\{A, G\}$  can be discovered simply by looking at the transitions from the pair of states on either 0 or on 1, and observing that (for one of those inputs) one state goes to  $C$  and the other does not. We can show  $\{A, G\}$  is distinguishable on the next round, since on input 1 they go to  $F$  and  $E$ , respectively, and we already established that the pair  $\{E, F\}$  is distinguishable.

However, then we can discover no more distinguishable pairs. The three remaining pairs, which are therefore equivalent pairs, are  $\{A, E\}$ ,  $\{B, H\}$ , and  $\{D, F\}$ . For example, consider why we can not infer that  $\{A, E\}$  is a distinguishable pair. On input 0,  $A$  and  $E$  go to  $B$  and  $H$ , respectively, and  $\{B, H\}$  has not yet been shown distinguishable. On input 1,  $A$  and  $E$  both go to  $F$ , so there is no hope of distinguishing them that way. The other two pairs,  $\{B, H\}$  and  $\{D, F\}$  will never be distinguished because they each have identical transitions on 0 and identical transitions on 1. Thus, the table-filling algorithm stops with the table as shown in Fig. 4.9, which is the correct determination of equivalent and distinguishable states.